

COMBINATORICS OF GIT GRAPHS



Julien COURTIEL (Université de

Caen Normandie

WORK 1

with Paul DORBEC Romain LECOQ (Université de Gen Normandie)

work 2

with Martin PEPIN (Université de Goen Normandie)



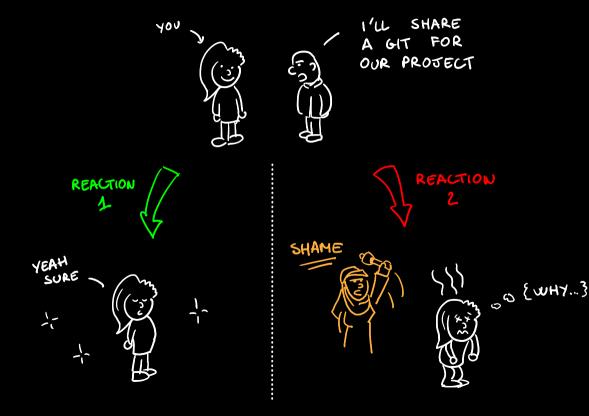
Git de France

ALEA Days 2025

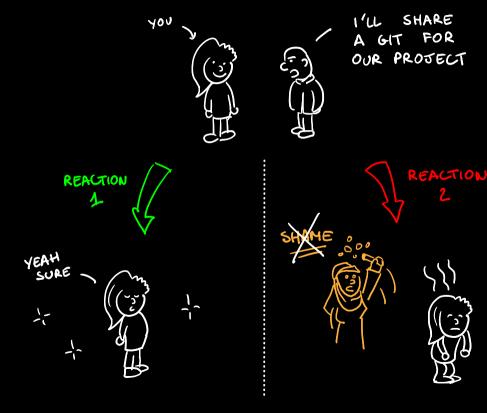
INTRODUCTION

How GIT WORKS

SITUATION PLAY



SITUATION PLAY









Situation play 2

You write an article with an obscure coauthor





REMOTE SERVER

COAUTHOR

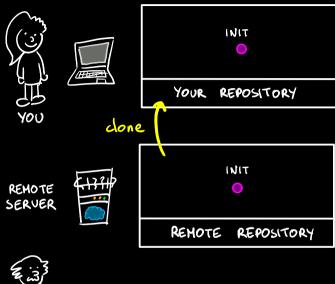


INIT



REMOTE REPOSITORY







COAUTHOR



INIT THEO

YOUR REPOSITORY

your article.pdf

REMOTE SERVER



REMOTE REPOSITORY

Toulouse has the best French accent

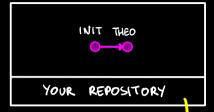
You Obscure Coauthor

Theorem 1. The Toulouse accent is the supremum of the set of all French accents.



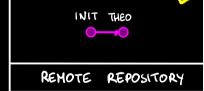
push





your article.pdf





Toulouse has the best French accent

You Obscure Coauthor

 ${\bf Theorem~1.}\ The\ Toulouse\ accent\ is\ the\ supremum\ of\ the\ set\ of\ all\ French\ accents.$





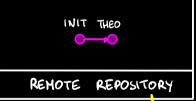
INIT THEO

YOUR REPOSITORY

your article.pdf

remote Server





REMOTE REPOSITORY

INIT THEO

OBSCURE'S REPOSITORY

Toulouse has the best French accent

You Obscure Coauthor

Theorem 1. The Toulouse accent is the supremum of the set of all French accents.





INIT THEO

TOUR REPOSITORY

your article.pdf

REMOTE SERVER

OBSCURE COAUTHOR



INIT THEO DEF

REMOTE REPOSITORY

push

typing detinition

OBSCURE'S REPOSITORY

Toulouse has the best French accent

You Obscure Coauthor

 ${\bf Theorem~1.~} \textit{The Toulouse accent is the supremum of the set of all French accents.}$



YOUR REPOSITORY

your article.pdf

REMOTE SERVER



REMOTE REPOSITORY In this document, A denotes the set of all French accents

You

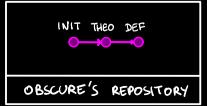
Toulouse has the best French accent

Obscure Coauthor

Definition 1. Define a relation order > on A such that for any two accents $a_1, a_2 \in A$, we have $a_1 > a_2$ if and only if accent a_1 is considered "greater than" accent as based on an objective set of criteria (including musicality, clarity, cultural significance).

Theorem 2. The Toulouse accent is the supremum of A.









YOUR REPOSITORY

REMOTE SERVER

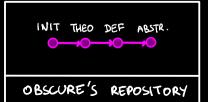
OBSCURE COAUTHOR





REMOTE REPOSITORY





your article.pdf

Toulouse has the best French accent

You Obscure Coauthor

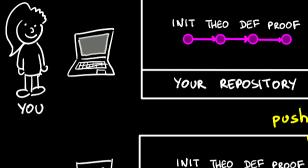
In this document, A denotes the set of all French accents.

Definition 1. Define a relation order > on \mathcal{A} such that for any two accents $a_1, a_2 \in \mathcal{A}$, we have $a_1 > a_2$ if and only if accent a_1 is considered "greater than" accent a_2 based on an objective set of criteria (including musicality, clarity, cultural significance).

Theorem 2. The Toulouse accent is the supremum of A.

Proof. To prove this, we will enumerate the various French accents in ${\cal A}$ and demonstrate that none can surpass the Toulouse accent.

- Paris Accent. Known for its elegance and sophistication, it's a bit too snobbish. Verdict: Not as warm as Toulouse.
- Marseille Accent. Full of passion and energy, this accent is like a nursery rhyme in heavy metal. Verdict: Too noisy.
- Alsace Accent. Although the appeal of old German is a little more fashionable these days, this accent is like a fusion restaurant that forgot the recipe. Verdict: Too confusing.
- Normand Accent. A bit rustic, this accent has a certain charm, but it often sounds like it's still trying to figure out where it parked its tractor. Verdict: Not chic.



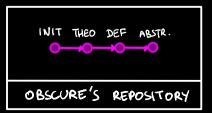
REMOTE SERVER

COAUTHOR









REPOSITORY

REMOTE

your article.pdf

Toulouse has the best French accent

Obscure Coauthor You

In this document. A denotes the set of all French accents.

Definition 1. Define a relation order > on A such that for any two accents $a_1, a_2 \in A$, we have $a_1 > a_2$ if and only if accent a1 is considered "greater than" accent a2 based on an objective set of criteria (including musicality, clarity, cultural significance).

Theorem 2. The Toulouse accent is the supremum of A.

Proof. To prove this, we will enumerate the various French accents in A and demonstrate that none can surpass the Toulouse accent.

- 1. Paris Accent. Known for its elegance and sophistication, it's a bit too snobbish. Verdict: Not as warm as Toulouse.
- Marseille Accent. Full of passion and energy, this accent is like a nursery rhyme in heavy metal. Verdict: Too noisy.
- 3. Alsace Accent. Although the appeal of old German is a little more fashionable these days, this accent is like a fusion restaurant that forgot the recipe. Verdict: Too confusing.
- 4. Normand Accent. A bit rustic, this accent has a certain charm, but it often sounds like it's still trying to figure out where it parked its tractor. Verdict: Not chic.

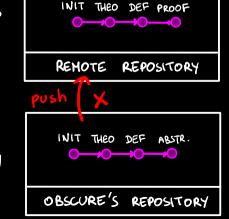




remote Server

OBSCURE COAUTHOR





your article.pdf

Toulouse has the best French accent

You Obscure Coauthor

In this document, A denotes the set of all French accents.

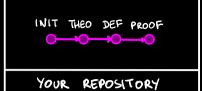
Definition 1. Define a relation order > on \mathcal{A} such that for any two accents $a_1, a_2 \in \mathcal{A}$, we have $a_1 > a_2$ if and only if accent a_1 is considered "greater than" accent a_2 based on an objective set of criteria (including musicality, clarity, cultural significance).

Theorem 2. The Toulouse accent is the supremum of \mathcal{A} .

Proof. To prove this, we will enumerate the various French accents in ${\cal A}$ and demonstrate that none can surpass the Toulouse accent.

- Paris Accent. Known for its elegance and sophistication, it's a bit too snobbish. Verdict: Not as warm as Toulouse.
- Marseille Accent. Full of passion and energy, this accent is like a nursery rhyme in heavy metal. Verdict: Too noisy.
- Alsace Accent. Although the appeal of old German is a little more fashionable these days, this accent is like a fusion restaurant that forgot the recipe. Verdict: Too confusing.
- Normand Accent. A bit rustic, this accent has a certain charm, but it often sounds like it's still trying to figure out where it parked its tractor. Verdict: Not chic.

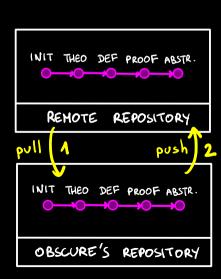




SERVER OBSCURE

REMOTE

COAUTHOR



your article.pdf

Toulouse has the best French accent

You Obscure Coauthor

In this document, A denotes the set of all French accents.

Definition 1. Define a relation order > on A such that for any two accents $a_1, a_2 \in A$, we have $a_1 > a_2$ if and only if accent a_1 is considered "greater than" accent a_2 based on an objective set of criteria (including musicality, clarity, cultural significance).

Theorem 2. The Toulouse accent is the supremum of \mathcal{A} .

Proof. To prove this, we will enumerate the various French accents in ${\cal A}$ and demonstrate that none can surpass the Toulouse accent.

- Paris Accent. Known for its elegance and sophistication, it's a bit too snobbish. Verdict: Not as warm as Toulouse.
- Marseille Accent. Full of passion and energy, this accent is like a nursery rhyme in heavy metal. Verdict: Too noisy.
- Alsace Accent. Although the appeal of old German is a little more fashionable these days, this accent is like a fusion restaurant that forgot the recipe. Verdict: Too confusing.
- Normand Accent. A bit rustic, this accent has a certain charm, but it often sounds like it's still trying to figure out where it parked its tractor. Verdict: Not chic.





your article.pdf

Your repository

poll

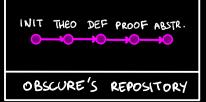
REMOTE SERVER





REMOTE REPOSITORY





Toulouse has the best French accent

You Obscure Coauthor

Abstract

This article demonstrates that the Toulouse accent is the suprenum of the set of French accents A by exhaustively enumerating its competitors, ultimately revealing that none can match its warmth, charm, and delightful character.

In this document, A denotes the set of all French accents.

Definition 1. Define a relation order > on A such that for any two accents $a_1, a_2 \in A$, we have $a_1 > a_2$ if and only if accent a_1 is considered "greater than" accent a_2 based on an objective set of criteria (including musicality, clarity, cultural significance).

Theorem 2. The Toulouse accent is the supremum of A.

Proof. To prove this, we will enumerate the various French accents in \mathcal{A} and demonstrate that none can surpass the Toulouse accent.

- Paris Accent. Known for its elegance and sophistication, it's a bit too snobbish. Verdict: Not as warm as Toulouse.
- Marseille Accent. Full of passion and energy, this accent is like a nursery rhyme in heavy metal. Verdict: Too noisy.





YOUR REPOSITORY

REMOTE SERVER

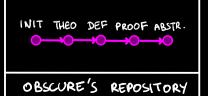




REMOTE REPOSITORY



COAUTHOR



your article.pdf

Toulouse has the best French accent

You Obscure Coauthor

Abstract

This article demonstrates that the Toulouse accent is the suprenum of the set of French accents A by exhaustively enumerating its competitors, ultimately revealing that none can match its warmth, charm, and delightful character.

In this document, A denotes the set of all French accents.

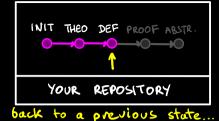
Definition 1. Define a relation order > on A such that for any two accents $a_1, a_2 \in A$, we have $a_1 > a_2$ if and only if accent a_1 is considered "greater than" accent a_2 based on an objective set of criteria (including musicality, clarity, cultural significance).

Theorem 2. The Toulouse accent is the supremum of A.

Proof. To prove this, we will enumerate the various French accents in \mathcal{A} and demonstrate that none can surpass the Toulouse accent.

- Paris Accent. Known for its elegance and sophistication, it's a bit too snobbish. Verdict: Not as warm as Toulouse.
- Marseille Accent. Full of passion and energy, this accent is like a nursery rhyme in heavy metal. Verdict: Too noisy.





your article.pdf

Toulouse has the best French accent

You Obscure Coauthor

In this document, A denotes the set of all French accents.

Definition 1. Define a relation order > on A such that for any two accents $a_1, a_2 \in A$, we have $a_1 > a_2$ if and only if accent a_1 is considered "greater than" accent a_2 based on an objective so of criteria (including musicality, clarity, cultural significance).

Theorem 2. The Toulouse accent is the supremum of A.

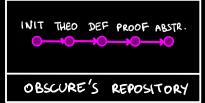
remote Server



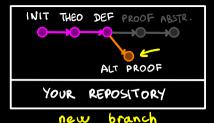




OBSCURE COAUTHOR





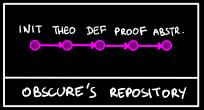


REMOTE SERVER

INIT THEO DEF PROOF ABSTR.

REMOTE REPOSITORY





your article.pdf

Toulouse has the best French accent

You Obscure Coauthor

In this document, A denotes the set of all French accents.

Definition 1. Define a relation order > on A such that for any two accents $a_1, a_2 \in A$, we have $a_1 > a_2$ if and only if accent a_1 is considered "greater than" accent a_2 based on an objective set of criteria (including musicality, clarity, cultural significance).

Theorem 2. The Toulouse accent is the supremum of \mathcal{A} .

Proof. (shorter proof??)

Assume, for the sake of contradiction, that the Toulouse accent is **not** the supremum of \mathcal{A} . Then, there exists an accent $a \in \mathcal{A}$ such that a is greater than the Toulouse accent.

However, upon hearing the Toulouse accent, any listener is irresistibly charmed and cannot help but prefer it over a. This contradicts our assumption that a is greater.

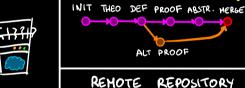
Thus, we conclude that our initial assumption is false, and the Toulouse accent must indeed be the supremum of A.





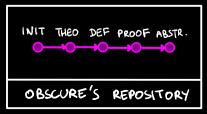
YOUR REPOSITORY

REMOTE SERVER





COAUTHOR



your article.pdf

Toulouse has the best French accent

Von Obscure Coauthor

Abstract

This article demonstrates that the Toulouse accent is the supremum of the set of French accents A by exhaustively enumerating its competitors, ultimately revealing that none can match its warmth, charm, and delightful character.

In this document. 4 denotes the set of all French accents.

Definition 1. Define a relation order > on A such that for any two accents $a_1, a_2 \in A$, we have $a_1 > a_2$ if and only if accent a1 is considered "greater than" accent a2 based on an objective set of criteria (including musicality, clarity, cultural significance).

Theorem 2. The Toulouse accent is the supremum of A.

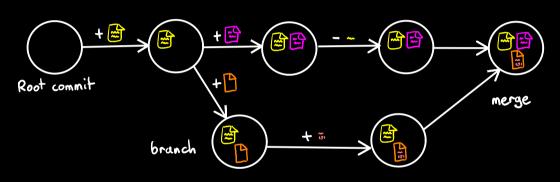
Proof. Assume, for the sake of contradiction, that the Toulouse accent is **not** the supremum of A. Then, there exists an accent $a \in A$ such that a is greater than the Toulouse accent.

However, upon hearing the Toulouse accent, any listener is irresistibly charmed and cannot help but prefer it over a. This contradicts our assumption that a is greater.

Thus, we conclude that our initial assumption is false, and the Toulouse accent must indeed be the supremum of A.

IN SUMMARY ...

◆ git is a Version Control System (VCS):
it stores all the project states over time.



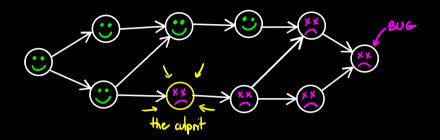
Git lets you create parallel development branches that can be integrated later.

This forms a Directed Acyclic Graph (DAG), where the vertices are the project states, also named commits.

PART II GIT BISECT



PROBLEM: FINDING A REGRESSION



Input A commit graph in which a commit is known to be bugged, the other commits are bugged or clean (= bug-free)

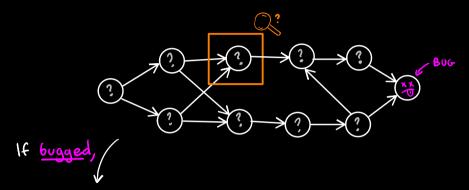
Question Which commit has originally introduced the bug?

Assumptions . If a parent of a commit is bugged, then the commit is bugged.

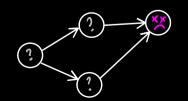
· Only one commit has introduced the bug, namely the faulty commit (or regression)

HOW TO INVESTIGATE

Unique operation: QUERY of a commit with unknown status



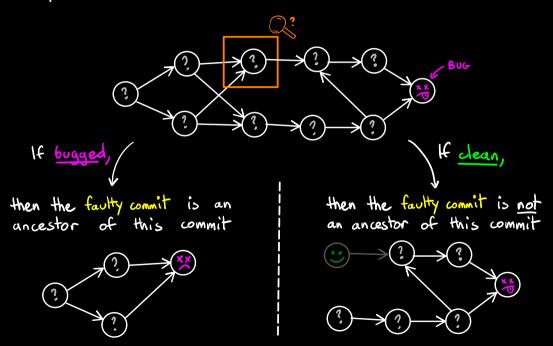
then the faulty commit is an ancestor of this commit



ancestor of a vertex v = v or an ancestor of a parent of v

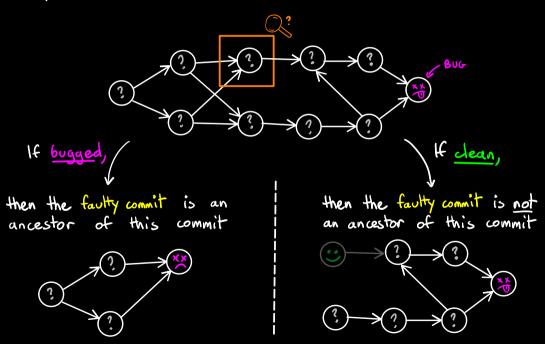
HOW TO INVESTIGATE

Unique operation: QUERY of a commit with unknown status



HOW TO INVESTIGATE

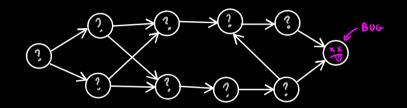
Unique operation: QUERY of a commit with unknown status



The faulty commit is found whenever there remains only I suspect.

REGRESSION SEARCH PROBLEM

Input: a DAG where each vertex has an unknown status, except one, which is bugged.



Output: A strategy that finds the faulty commit with a minimal number of queries in the worst-case scenario = optimal strategy

(the faulty commit can be any ancestor of the bugged vertex)

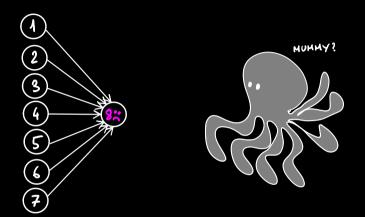
In real life, queries are costly.

FIRST EXAMPLE : CHAINS

optimal strategy = binary search

More generally, number of queries in an optimal strategy in a chain of length $n = \lceil \log_2(n) \rceil$

SECOND EXAMPLE : OCTOPUSES

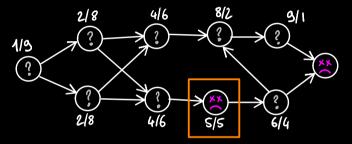


optimal strategy = whatever

More generally, number of queries in an optimal strategy in an octopus of size n = n-1



STEP 1: Compute the number of ancestors/non-ancestors for each commit

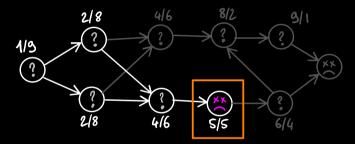


STEP 2: Query the vertex with the most balanced ratio

STEP 3:



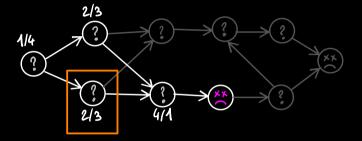
STEP 1: Compute the number of ancestors/non-ancestors for each commit



STEP 2: Query the vertex with the most balanced ratio



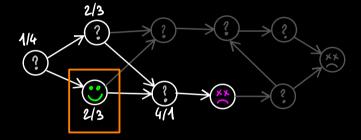
STEP 1: Compute the number of ancestors/non-ancestors for each commit



STEP 2: Query the vertex with the most balanced ratio



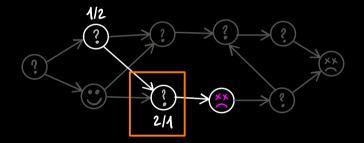
STEP 1: Compute the number of ancestors/non-ancestors for each commit



STEP 2: Query the vertex with the most balanced ratio



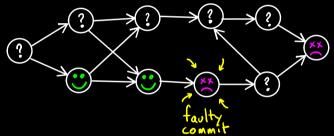
STEP 1: Compute the number of ancestors/non-ancestors for each commit



STEP 2: Query the vertex with the most balanced ratio



STEP 1: Compute the number of ancestors/non-ancestors for each commit



STEP 2: Query the vertex with the most balanced ratio

STEP 3: Delete the innocent commits and recurse.

QUESTION: Does git bisect always give an optimal strategy?

HOW GOOD IS GIT BISECT?

QUESTION: Does git bisect always give an optimal strategy?

[NO] The Regression Search Problem is NP-complete.

[Carmo Donadelli Kohayakawa Laber 2004] [We've also proved it!]

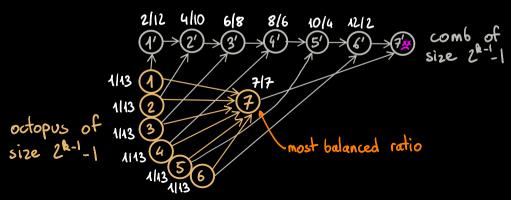
But is it really that bad? YEAH

Proposition For any R, there exists a DAG such that an optimal strategy uses R queries and git bisect always uses 2^{k-1}-1 queries.

HOW GOOD IS GIT BISECT?

Proposition For any R, there exists a DAG such that an optimal strategy uses R queries and git bisect always uses 2^{k-1}-1 queries.

Proof for k=4:



BACK TO REALITY?

Octopus substructures are unrealistic

Classical git graph:	
Main O YI.I Develop	V2.0
Feature 2	O c -

Usually, we never merge more than 2 branches.

(Otherwise it is called an octopus merge

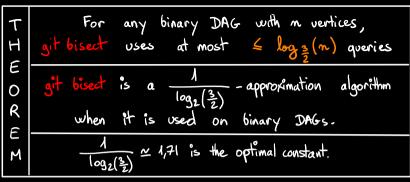
BINARY DAGS

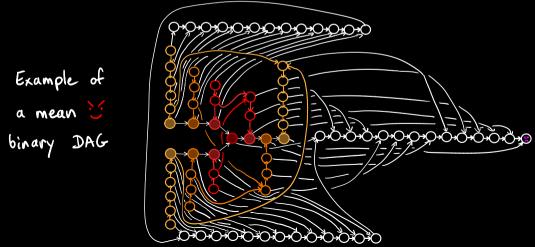
Definition binary DAG = DAG where the vertices have indegree ≤ 2

Ex: O Bud

H I I	
	git bisect is a $\frac{1}{\log_2(\frac{3}{2})}$ -approximation algorithm when it is used on binary DAGs.
EM	$\frac{1}{\log_2(\frac{3}{2})} \simeq 1.71$ is the optimal constant.

BINARY DAGS





GOLDEN BISECT

We've improved git bisect in the worst-case scenario

-> new algorithm: \$ golden bisect \$

Idea behind of therwise, query some special vertex even if it has not best score.

$$\phi$$
 = golden ratio = $(1+\sqrt{5})/2$

THEORE	For any binary DAG with n vertices, golden bisect uses at most $\leq \log_{\frac{1}{2}}(n)+1$ queries
	golden bisect is a $\frac{1}{\log_2(\phi)}$ -approximation algorithm when it is used on binary DAGs.
M	$\frac{1}{\log_2(\phi)} \simeq 1,44$ is the optimal constant.

WHERE IS MY ALEA?

Number of queries in the average-case scenario?

- · DAG: binary, fixed, n vertices (not random)
- <u>Position of the</u> Uniformly at random faulty commit: amongst vertices of the DAG

- NEW -- RESULT: [C., Dorbec, Jugé]

THEOREM

- . The average number of queries for git bisect is $\leq \log_{\frac{3}{2}} 2^{V_3} (n)$
- . The average number of queries for golden bisect is $\leq \log_2 \phi^{1+\phi^{-2}}(n)$

QUESTION

Have we really proved that git bisect is bad?

- -> The examples of DAGs where git bisect performs poorly shouldn't occur in real life
- -> Sometimes git bisect is better than golden bisect

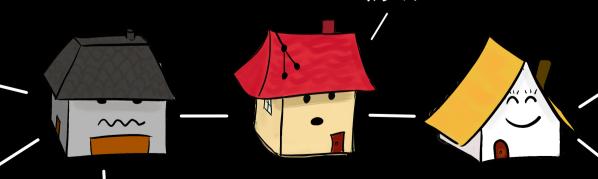
More relevant (?) What is the average-case question complexity of git bisect?

This question calls for many more, notably what is a random Git graph?

PART III RANDOM GIT GRAPHS

ongoing work

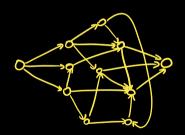
IT'S MY NEIGHBORHOOD!

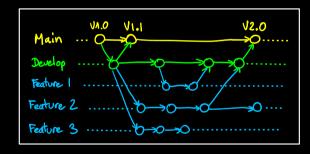


WHICH GRAPHS TO CONSIDER?

In ogit, every DAG without restriction can be generated ...

... but many projects follow a workflow





In the following, we consider a simple workflow but widely used in industry: the feature branch workflow

GIT GRAPH

DEFINITION

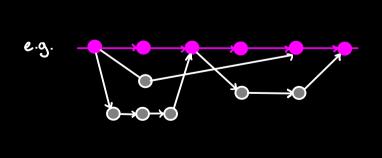
feature branch)

· a main branch (path of magenta vertices)

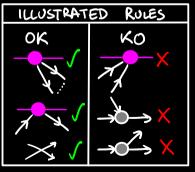
•0,1 or several feature branches, paths of >1 white vertices starting and ending on magenta vertices

indegree ≤ 2 for all vertices

previously defined in [Lecoq 2024]



= DAG with



GIT GRAPH

DEFINITION

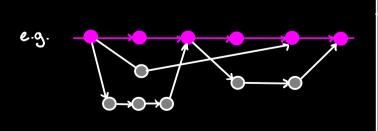
(feature branch)
Git graph

a main branch (path of magenta vertices)

. O, 1 or several feature branches, paths of >1 white vertices starting and ending on magenta vertices

indegree ≤ 2 for all vertices

previously defined in [Lecoq 2024]

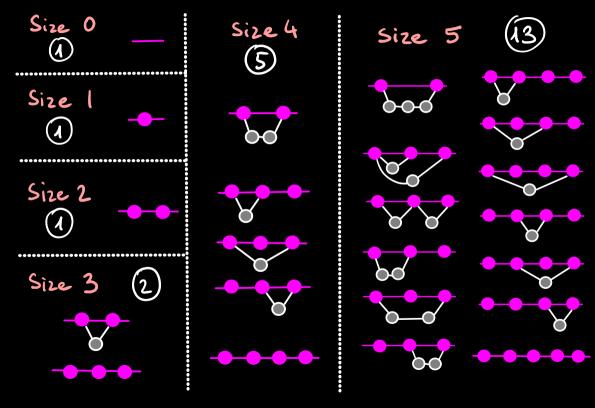


= DAG with

GOALS

- · Counting Git graphs (exactly or asymptotically)
 - · Random generation of Git graphs

ALL SHALL GIT GRAPHS



RECURSIVE DECOMPOSITION

Decomposition

 $G(r_{3},u) = 1 + r_{3}u G(r_{3},u) + \frac{r_{3}^{2}u^{2}}{1-r_{3}} \frac{\partial G}{\partial u} (r_{3},u)$ where $G(r_{3},u) = \sum_{n \geq 0} \sum_{k \geq 0} r_{3}n_{k} r_{3}^{n}u_{k}$

THREE ANGLES OF ATTACK

Sandwich method

Finding subsets and supersets of Git graphs easier to count

Counting Gaphs

Ennoblement method

making the generating function less ordinary

Freezer method



Counting by (temporarily)
fixing the number of
magenta vertices

ENNOBLEMENT METHOD

 $g_{n,k} = g_{n-1,k-1} + \sum_{l \geq 1} (k-1) g_{n-1-l,k-1}$ Recurrence-Differential $G(z,u) = 1 + zu G(z,u) + \frac{z^2 u^2}{1-z} \frac{\partial G}{\partial u}(z,u)$ Equation Usual trick: Borel transform Ordinary

2, 80 ml 2 ml m, k zo gm, k 22 uk Generating Generating Function Function (-(m, gr)) not analytic X Bore

on u

analytic V

onalytic, but no pretty equation X Differential Equation for G $G(z_1,u) = \sum_{m,k>0} \frac{g_{m,k}}{k!} z_m^{m,k}$ 3G = 3G + 32 M 3G

$$(z,u) = 1 + zu G(z,u)$$

Borel transform

Exponential

ENNOBLEMENT METHOD

Recurrence
$$3n, k = 9n-1, k-1 + \sum_{l \ge 1} (k-1) g_m - 1 - l, k-1$$

Differential $G(x_j, u) = 1 + x_j u G(x_j, u) + \frac{x^2 u^2}{1 - x_j} \frac{3G}{3u} (x_j, u)$

$$\widetilde{G}(z_{3},u) = \sum_{n,k>0} \frac{g_{n,k}}{k!} \frac{g_{n,k}}{g_{n,k}} \frac{g$$

DIFFERENT PERSPECTIVE

Definition

set of cycles of

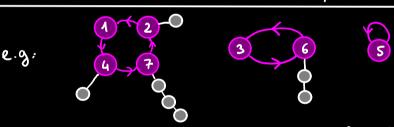
magenta vertices labeled from 1 to h

where a chain of white unlabeled vertices

is attached to each magenta vertex,

except to the ones having the

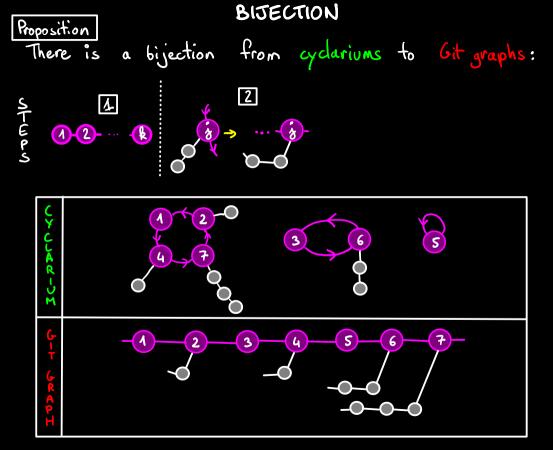
smallest label in their cycles.

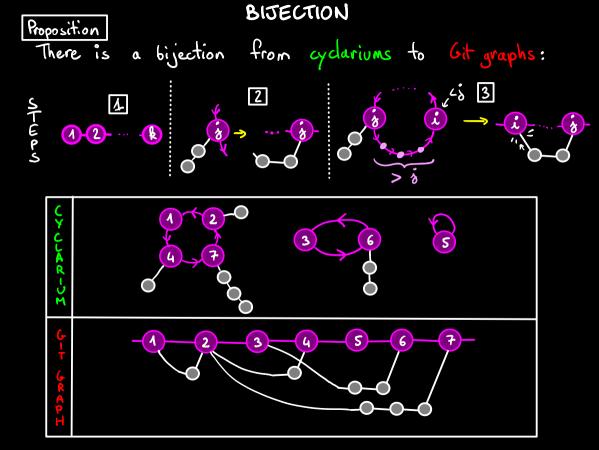


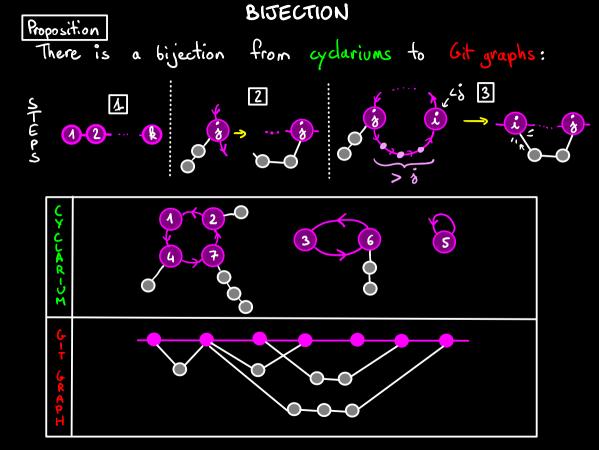
Is there a combinatorial explanation for the formula

$$\widetilde{G}(r_{3/N}) = \left(1 - \frac{r_{3}^{2} n}{1 - r_{3}}\right)^{-\frac{1 - r_{3}}{r_{3}}} = \exp\left(\frac{1}{\frac{r_{3}}{1 - r_{3}}} \ln\left(\frac{1}{1 - n_{3} \frac{r_{3}}{r_{3}}}\right)\right)$$

It's the generating function of cyclariums!







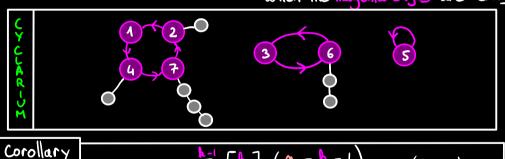
BIJECTION

Proposition

There is a bijection from cyclariums to Git graphs sending vertices —> vertices

magenta vertices —> magenta vertices

cycles —> free vertices of indegree \(\) i.e magenta vertices of indegree \(\) cycle lengths —> magenta vertices in the connected components when the magenta edges are erased.



orollary

$$g_{m,k} = \sum_{f=1}^{k-1} \begin{bmatrix} k \\ f \end{bmatrix} \begin{pmatrix} n-k-1 \\ k-f-1 \end{pmatrix}$$

(kg_{m,k} = n_{m} number of Git graphs counted by vertices & magenta and $\begin{bmatrix} \vdots \end{bmatrix} = (u_{n}signed)$ Stirling number of 1st kind

RANDOM MODEL

"Boltzmann model" (exponential in u, ordinary in 2)

We wish to draw a Git graph 8 with a weight proportional to represent to white wertices in 8 the magenta vertices in 8 (# magenta vertices in 8)!

where $G(z,u) = \sum_{m,k \geq 0} \frac{2^{m,k}}{k!} z^{mu} = \left(1 - \frac{3^2 u}{1 - 3}\right)^{-\frac{3^2 u}{1 - 3}}$

Examples
$$P(-) = \frac{1}{\tilde{G}(y_0 u)} P(-) = \frac{3}{\tilde{G}(y_0 u)} P(-) = \frac{5}{\tilde{G}(y_0 u)} P(-) = \frac{5}{\tilde{G}(y_0 u)} \frac{1}{\tilde{G}(y_0 u)} \frac{1}{\tilde{G}(y_0 u)} \frac{1}{\tilde{G}(y_0 u)} P(-) = \frac{5}{\tilde{G}(y_0 u)} \frac{1}{\tilde{G}(y_0 u)}$$

*: in the disk of convergence of G

RANDOM MODEL

Proposition Let 8 be a random Git Graph sampled with respect to the previous Boltzmann model, conditioned to have size n $\mathbb{E}(\# \text{ magenta vertices}(8)) \sim \frac{1-q_u}{2-q_u}$ V (# magenta vertices (8)) ~ Pu(1-Pu) n where $Qu = \frac{(2-Pu)^3}{2u}$ Proof: Transfer Theorem from $\widetilde{G}(r_{0,1}u) = \left(1 - \frac{3^2u}{r_0}\right)^{\frac{1-r_0}{r_0}}$

Consequence: A random generator for Git graphs with
$$\approx$$
 n vertices and \propto k magenta vertices ($k \le \frac{n}{2}$)

1. Tune u so that $\frac{1-\varrho_u}{2-\varrho_u} = \frac{k}{n}$ 3. Make a Boltzmann samplen with parameters $\frac{1}{2}$ and $\frac{1}{2}$.

The so that $\frac{1-\varrho_u}{2-\varrho_u} = \frac{1-\varrho_u}{n}$ for cyclariums.

2. Tune 2 so that $z = \rho_u - \frac{1 - \rho_u}{n}$ 4. Bijection to Git graphs

THREE ANGLES OF ATTACK

Sandwich method

Finding subsets and supersets of Git graphs easier to count

Counting Gat Graphs

Ennoblement method

making the generating function less ordinary

Freezer method *

Counting by (temporarily)
fixing the number of
magenta vertices

IT'S EASIER WHEN YOU FIX &

$$G_k(z) := \sum_{n \geqslant 0} g_{n,k} z_n^n$$

= Generating Function of Git graphs where the number & of magenta vertices is fixed

Claim:
$$G_{k}(y) = y_{k} \prod_{s=0}^{k-1} (1 + \frac{y_{k}}{3})$$

$$= \frac{y_{k}}{(1-y_{k})^{k}} \frac{\Gamma(k+y_{k}/(1-y_{k}))}{\Gamma(y_{k}/(1-y_{k}))}$$

Cauchy integral $g_{n,k} = \frac{1}{2\pi i} \int \frac{2k-n-1}{(1-2g)^k} \frac{\Gamma(k+3/(1-3g))}{\Gamma(3/(1-3g))} dy$

-> invitation to saddle-point ...

ASYMPTOTIC ESTIMATE

$$g_{n,k} = \frac{1}{2\pi i} \int \frac{3^{2k-m-1}}{(1-3)^k} \frac{\Gamma(k+3/(1-3))}{\Gamma(3/(1-3))} ds$$

NEW PRESULT.

theorem The number of Git graphs with n vertices is asymptotically equivalent to

$$g_{n} \sim \frac{e^{1/8}}{2} \left(\frac{m}{2e}\right)^{\frac{n}{2}} \exp\left(\frac{1}{2} \ln\left(\frac{n}{2}\right) \sqrt{\frac{n}{2}} + \sqrt{2n} + \frac{\left(\ln\frac{n}{2}\right)^{2}}{32}\right) \left(\frac{n}{2}\right)^{-\frac{5}{8}}$$

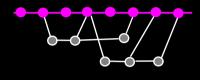
gn, k is maximal for $k \approx \frac{n}{2} + \frac{1}{4} ln \left(\frac{n}{2}\right) \sqrt{\frac{n}{2}}$ and we should have a Gaussian local limit law ...

PERSPECTIVES ABOUT RANDOM GIT GRAPHS

- -> Other random models
 - · Other graph models

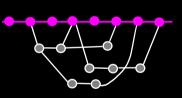
D Collaboration with Clement + Maréchal

Phoenix graphs



merged branches can be reborn

Fork Anywhere graphs



branches can be born anywhere but must be merged into main

More involved workflows



· more probabilistic models

LINKS BETWEEN GIT GRAPHS & GIT BISECT

Still to do: Average-case complexity of git bisect where the input is taken w.r.t the Boltzmann distribution.

But also: - Is there a polynomial algorithm for the Regression Problem when the input is a Git graph?

(We proved that git bisect fails to be optimal for some Git graphs)

- Is the Regression Problem NP-complete when the input RESULT: [Bulteau +]: No (if ETH is true) is binary?
- Other algorithms from Version Control Systems to be analyzed?

THANK YOU!

